Bishop operators and Diophantine approximation

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Short abstract. One of the oldest unsolved problems in functional analysis is the invariant subspace problem. E. Bishop proposed a family of potential counterexamples depending on an irrational parameter and 20 years later A. M. Davie strongly contradicted Bishop's intuition proving that it only could be a counterexample for Liouville numbers. The purpose of this talk is to illustrate this noticeable interplay between number theory and functional analysis and to present a recent joint work with E. Gallardo, M. Monsalve and A. Ubis.

Warning. The emphasis here is in the *application* of number theory not in the strength of the number theoretical results proven with this purpose.

Does every bounded linear operator on a (complex, ∞ -dim., separable) Hilbert space have a nontrivial closed invariant subspace?

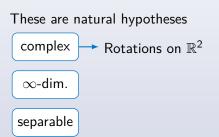
These are natural hypotheses



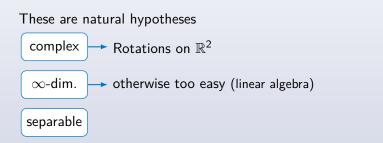




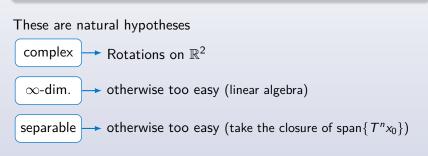
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Every known Hilbert space satisfying the hypotheses has an invariant subspace. Should we bet that invariant subspaces always exist?

• Yes? Lomonosov (1973). True for any operator commuting with a compact operator.

• No? Enflo (1976-1987), Read (1984). There are counterexamples in Banach spaces. In fact there exists an operator on ℓ^1 without invariant subspaces (Read, 1985).

Gossiping about the invariant subspace problem:

- Lomonosov's proof of his theorem is astonishingly simple and elegant. His paper is only two pages long.
- It made obsolete the previous partial results on the problem. Even during some years it was unclear if Lomonosov's proof had fully solved it because any known operator satisfied the hypotheses.
- Enflo's paper is 101 pages long, took 7 years between submission and publication and 10 years since the initial announcement.
- Read's first paper is 64 pages long and it was published while Enflo's paper was in the refereeing process.

• Yes? Lomonosov (1973). True for any operator commuting with a compact operator.

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The counterexamples are too complicated to give any insight. Let's ask for a third opinion

• Maybe (not)? Bishop (ca. 1950) suggested that the operator

 $f(t) \mapsto t f(\{t + \alpha\})$

on $L^2[0,1)$ with $\alpha \in \mathbb{R} - \mathbb{Q}$ and $\{\cdot\}$ = fractional part, could be a counterexample.

Bishop operators

Let us define the Bishop operators with a certain scaling

 $Tf(t) = et f(\{t + \alpha\})$ e = 2.718182...

It is the composition of two well-known operators: a shift operator (with a continuous spectrum) and a multiplication operator (not invertible and no eigenfunctions). If $\alpha \in \mathbb{Q}$ the subspace of functions with small support around $k\alpha$ is invariant.

Possible rationale for it:

The equidistribution of $\{k\alpha\}$ for $\alpha \notin \mathbb{Q}$ mixes the support and the multiplication avoid eigenfunctions. If this implies that $\{T^k f\}$ gives everything we have solved the problem in the negative.

Invari		

More than 20 years later, A.M. Davie dealt Bishop's intuition a heavy blow.

Theorem (Davie, 1974). If α is not a Liouville number then T has nontrivial invariant subspaces.

Nevertheless many authors have continued the research on T and its variants. It seems that it is a natural operator to consider.

The folklore conjecture is the opposite of the original one: T has nontrivial invariant subspaces for every α .

Gossiping about Davie's paper:

- Halmos (1974) was rather tough in Mathematical reviews. "Several canons of both the pedantic grammarian and the mathematical expositor are violated. There is no hint to what motivates the proof [...] the techniques make use of nontrivial facts about subjects as diverse as diophantine approximations, Banach algebras, and quasi-analytic classes. The proofs appear, however, to be correct [...]"
- ▶ Halmos (1987) praises Davie and the result in his photo book.
- Flattot (2008) says that the proof is "very elliptic".

The proof is based on a method introduced by Wermer and extended by Atzmon. When it is applied, Davie had to confront some basic questions about Diophantine approximation.

Atzmon theorem (weak version)

Theorem (Atzmon, 1984).
$$L: H \longrightarrow H$$
, $f, g \in H - \{0\}$ with
 $\|L^{\pm n}f\|, \|(L^*)^{\pm n}g\| \le W(n), \qquad W(n) = \exp\left(\frac{n}{(\log n)^{1+\epsilon}}\right)$

Atzmon theorem

+ technical condition (easy for Bishop's operators)

 \Rightarrow L has a nontrivial invariant subspace.

Sketch

Take $\phi(z) = \sum a_k z^k$, $\psi(z) = \sum b_k z^k$ with $\psi \cdot \phi = 0$ on S^1 and $a_k W(|k|)$, $b_k W(|k|)$ small (possible by Denjoy-Carleman Theorem).

Define $u = \phi(L)f$, $v = \psi(L^*)g$ ($\neq 0$ by the technical condition). $M = \text{span}\{L^n u\}$ is invariant and it is nontrivial because $v \in M^{\perp}$:

$$\langle L^n u, v \rangle = \langle \phi(L) L^n f, \psi(L^*) g \rangle = \langle (\psi \phi)(L) L^n f, g \rangle = 0.$$

Folklore conjecture. The Bishop operators $Tf(t) = et f(\{t + \alpha\})$ have nontrivial invariant subspaces for every $\alpha \notin \mathbb{Q}$.

Davie (1974). True for $|\alpha - \frac{a}{q}| > q^{-N}$, (N arbitrary).

..... (MacDonald, other authors)

Flattot (2008). True for
$$|\alpha - \frac{a}{q}| > \exp(-q^{1/3})$$
, $(\frac{1}{3} \leftrightarrow \frac{1}{2} - \epsilon)$.

Our contribution (joint with Gallardo, Monsalve, Ubis 2018).

- Short proof of Flattot result.
- True for $|\alpha \frac{a}{q}| > \exp(-q^{1-\epsilon})$.
- For $|\alpha \frac{a}{q}| \ge \exp\left(-\frac{Cq}{\log q}\right) \nexists f, g$ to which Atzmon theorem can be applied (no possibility of improvement with the known methods).

Diophantine approximation

$$Tf(t) = et f({t + \alpha})$$
 $e = 2.718182...$

A direct substitution proves

$$T^n f(t) = e^{L_n(t)} f(\lbrace t + n\alpha \rbrace), \quad T^{-n} f(\lbrace t + n\alpha \rbrace) = e^{-L_n(t)} f(t)$$

and similar formulas for T^* with

$$L_n(t) = \sum_{k=0}^{n-1} \left(1 + \log\{t + k\alpha\}\right) \quad \text{note } \int_0^1 \log = -1.$$

Idea (Davie). Choose the support of f in such a way that t is never close (mod. 1) to $m\alpha$, $m \in \mathbb{Z}$.

Bad guys. If α is very close to a rational, L_n is amplified.

Matching Flattot's result

Take
$$f = \mathbf{1}_{\mathcal{D}}$$
 with $\mathcal{D} = \{t : \langle t \pm n\alpha \rangle > \frac{1}{2019n^2}, n \in \mathbb{Z}^+\}$ (Davie)
 $a/q = \text{ convergent of } \alpha$
 $L_n(t) = \sum_{k=0}^{n-1} (1 + \log\{t + k\alpha\}) \ll q + \frac{n+q}{q} \log(n+q).$

Idea. Each *q*-block contributes $O(\log n + \log q)$.

Short proof of the result by Flattot

Atzmon theorem requires $L_n(t) \ll \frac{n}{(\log n)^{1+\epsilon}}$. It is achieved if $\frac{\log n}{q} \ll \frac{1}{(\log n)^{1+\epsilon}}$ for $n > q^{3/2}$. A/Q = next convergent, $q^{3/2} < n \le Q^{3/2} \rightsquigarrow \log Q \ll q^{1/2-\epsilon}$ and it is fulfilled for $|\alpha - \frac{a}{q}| > \exp(-q^{1/2-\epsilon})$.

Proofs

$$t\in \mathcal{D} \Rightarrow \langle t\pm nlpha
angle >\epsilon_{\it n} symp n^{-2}$$
 and essentially

$$L_n(t) \ll q + nq^{-1} \log (\epsilon_n^{-1} + q)$$
 is best possible.

Pessimistic view. Increase $\epsilon_n \Rightarrow |\mathcal{D}| = 0$ we ran out of points $\stackrel{(?)}{\Rightarrow}$ No chance of improvement.

Dreaming. With the nonsensical choice $\epsilon_n = 1$, Atzmon theorem imposes $\frac{n}{q} \log q \ll \frac{n}{(\log n)^{1+\epsilon}}$ and for $q^{3/2} < n \leq Q^{3/2}$ gives $\log Q \ll q^{1-\epsilon}$ (Our result).

Hope. The dream only requires $\log(n+q) \rightarrow \log q$ to come true.

$$L_n(t) = \sum_{k=0}^{n-1} (1 + \log\{t + k\alpha\}) \qquad \frac{a}{q}, \ \frac{A}{Q} \text{ conv. } \alpha \notin \mathbb{Q}$$
$$\mathcal{D} = \{t : \langle t \pm n\alpha \rangle \gg \frac{1}{n^2}, \ \forall n \in \mathbb{Z}^+\}$$

Dreamed bound. $L_n(t) \ll q + \frac{n}{q} \log q$ for $t \in \mathcal{D}$

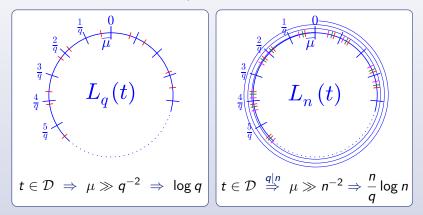
Proposition. For
$$Q \gg q^4$$
, $n \leq Q^{3/2}$, $t \in \mathcal{D}$

$$L_n(t) \ll q + rac{n'}{q}\log q + rac{n+Q}{Q}\log r$$

n' = remainder when n is divided by Q.

Idea of the proof

$$L_n(t) = \sum_{k=0}^{n-1} \left(1 + \log\{t + k\alpha\} \right)$$

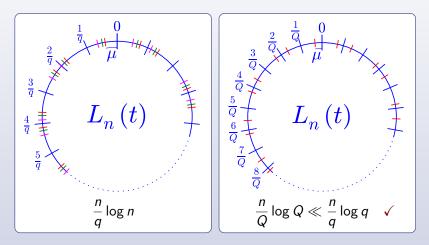


Target. Replace $\log n$ by $\log q$.

Range to keep in mind log
$$Q \ll q^{1-\epsilon}$$

$$(q^{3/2} < n < Q^{3/2})$$

Case $n \ge Q$ [Use next convergent]



Range to keep in mind log
$$Q \ll q^{1-\epsilon}$$

$$(q^{3/2} < n < Q^{3/2})$$

Case $n \leq Q$ [Perturbation] $\alpha = \frac{a}{q} + \frac{\delta}{qQ}$, n = qk + j

$$\{t+n\alpha\} = \{t+j\alpha+\frac{k\delta}{Q}\}, \qquad t \in \mathcal{D} \Rightarrow \{t+j\alpha\} \gg \frac{1}{q^2}$$

If $n/qQ \ll q^{-2}$ then $k/Q \ll q^{-2}
ightarrow$ perturbed version of $L_q(t)$

$$L_n(t) o rac{n}{q} \widetilde{L}_q(t).$$

Then essentially

$$n \ll rac{Q}{q} \quad \Rightarrow \quad L_n(t) \ll rac{n}{q} \log q \quad \checkmark$$

It remains to consider $n \gg Q/q$

Range to keep in mind log
$$Q \ll q^{1-\epsilon}$$

$$(q^{3/2} < n < Q^{3/2})$$

Case $Q/q \ll n \le Q$ [$j \ne 0 \rightarrow$ small scale] $\alpha = \frac{a}{q} + \frac{\delta}{qQ}$, n = qk + j

$$L_n(t)
ightarrow \sum_{k=0}^{n/q-1} \sum_{j=0}^{q-1} \Big(1 + \log \big(\mu + \frac{j}{q} + \frac{k\delta}{Q} \big) \Big), \qquad \log \mu \ll \log n.$$

If $j \neq 0$, k/Q is in a smaller scale than $j/q \Rightarrow$ no interference. Essentially

$$L_n(t) - (\text{contrib. } j = 0) \ll \log n + \frac{n}{q} \log q \ll \frac{n}{q} \log q \quad \checkmark$$

Range to keep in mind log
$$Q \ll q^{1-\epsilon}$$

$$(q^{3/2} < n < Q^{3/2})$$

$$\sum_{k=0}^{n/q-1} \left(1 + \log\left(\mu + \frac{k\delta}{Q}\right)\right)$$

with
$$\log \mu \ll \log n$$
. Then
(contrib. $j = 0$) $\ll \frac{n}{q} + \log n + \sum_{k=1}^{n/q-1} \log \frac{qQ}{n} + \sum_{k=1}^{n/q-1} \log \frac{n}{qk}$
 $\ll \frac{n}{q} + \log n + \text{(trivial)} + \text{(Stirling)}$
 $\ll \frac{n}{q} + \log n + \frac{n}{q} \log q + \frac{n}{q} \ll \frac{n}{q} \log q$.

Sharpness

$$\left| \alpha - \frac{a}{q} \right| \neq \exp \left(- C \frac{q}{\log q} \right) \Rightarrow \nexists f$$
 under Atzmon theorem.

The proof employs the result:

$$\frac{q^2\log q}{\epsilon^2} < n < \epsilon^2 \frac{Q}{q} \Rightarrow L_n(t - n\alpha), -L_n(t) > \epsilon \frac{n}{q}\log q$$

for $t \in [0,1)$ except for a set S of measure $O(\epsilon)$.

$$\|T^n f\| + \|T^{-n} f\| \gg \inf_{t \notin S} \left(e^{L_n(t-n\alpha)} + e^{-L_n(t)} \right) \|f\| \rightarrow \text{ too large.}$$

A copy of these slides is available in

https://www.uam.es/fernando.chamizo

Thank you for your attention!